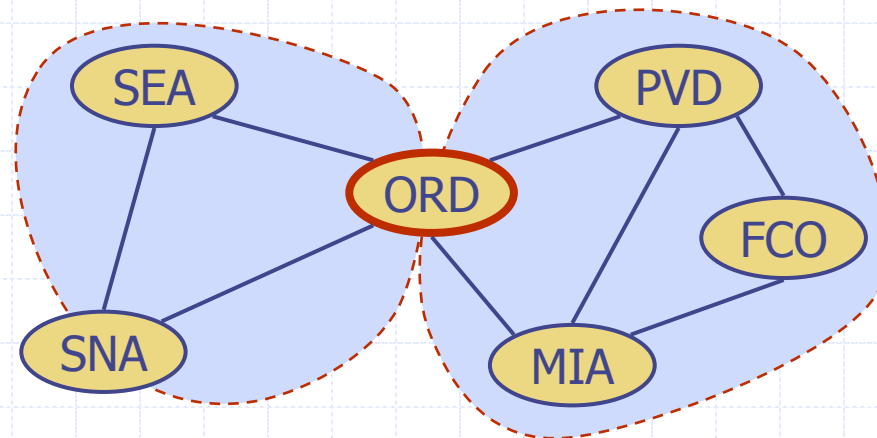


Biconnectivity



Outline and Reading

◆ Definitions (§6.3.2)

- Separation vertices and edges
- Biconnected graph
- Biconnected components
- Equivalence classes
- Linked edges and link components

◆ Algorithms (§6.3.2)

- Auxiliary graph
- Proxy graph

Separation Edges and Vertices

◆ Definitions

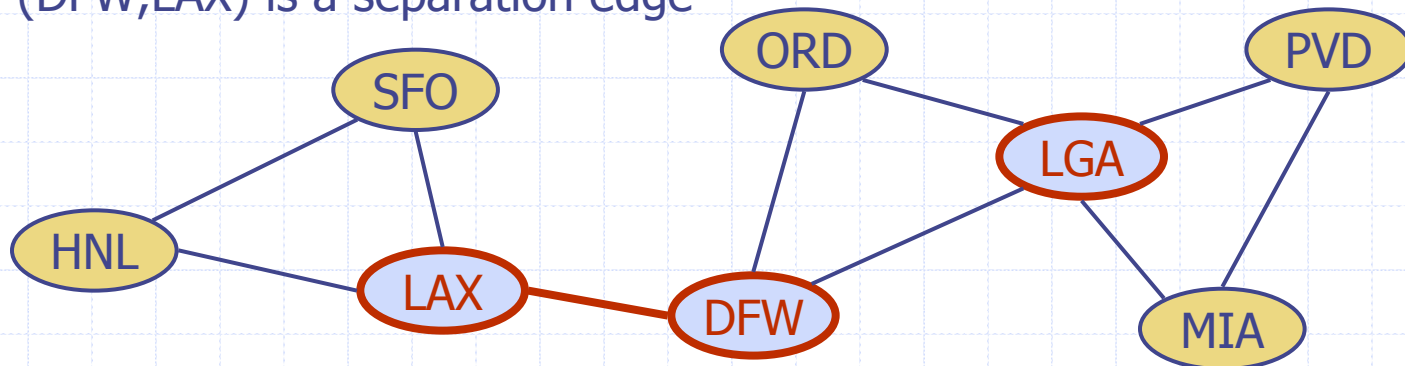
- Let G be a connected graph
- A separation edge of G is an edge whose removal disconnects G
- A separation vertex of G is a vertex whose removal disconnects G

◆ Applications

- Separation edges and vertices represent single points of failure in a network and are critical to the operation of the network

◆ Example

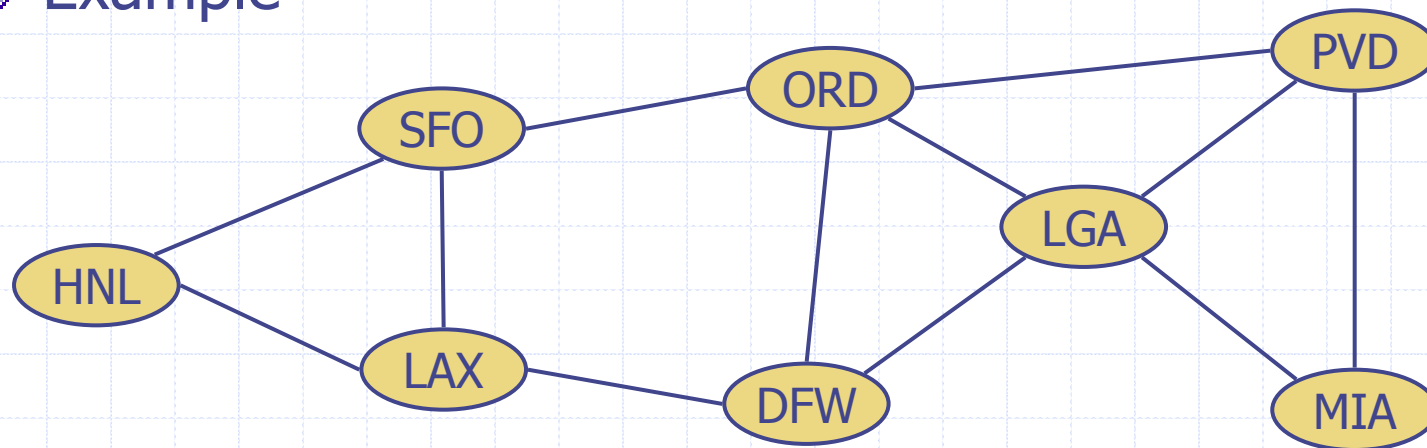
- DFW, LGA and LAX are separation vertices
- (DFW,LAX) is a separation edge



Biconnected Graph

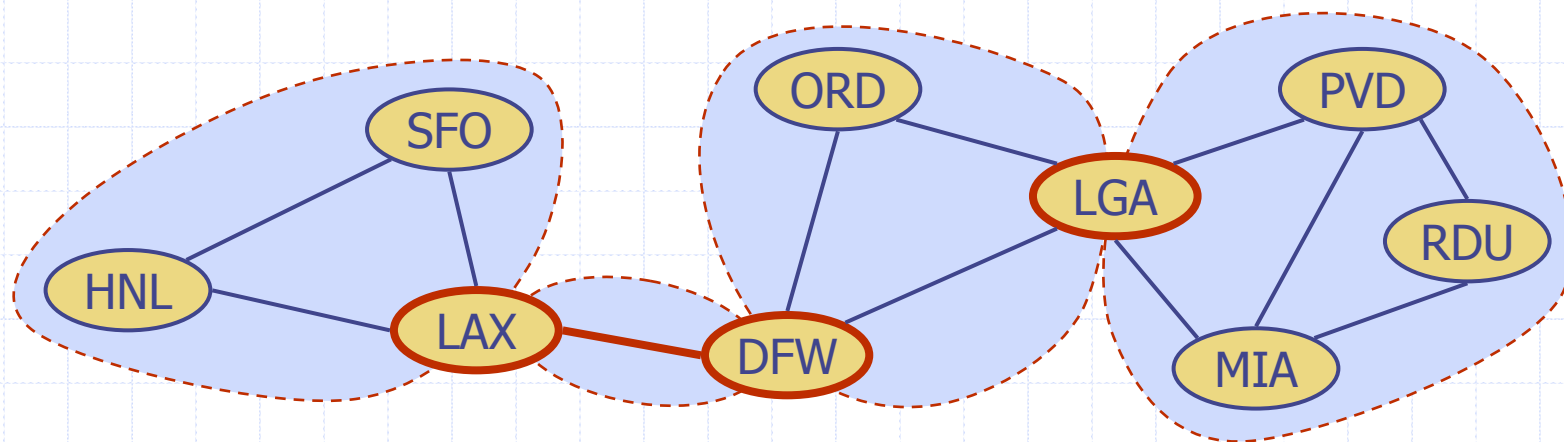
- ◆ Equivalent definitions of a biconnected graph G
 - Graph G has no separation edges and no separation vertices
 - For any two vertices u and v of G , there are two disjoint simple paths between u and v (i.e., two simple paths between u and v that share no other vertices or edges)
 - For any two vertices u and v of G , there is a simple cycle containing u and v

- ◆ Example



Biconnected Components

- ◆ Biconnected component of a graph G
 - A maximal biconnected subgraph of G , or
 - A subgraph consisting of a separation edge of G and its end vertices
- ◆ Interaction of biconnected components
 - An edge belongs to exactly one biconnected component
 - A nonseparation vertex belongs to exactly one biconnected component
 - A separation vertex belongs to two or more biconnected components
- ◆ Example of a graph with four biconnected components



Equivalence Classes

- ◆ Given a set S , a relation R on S is a set of ordered pairs of elements of S , i.e., R is a subset of $S \times S$
- ◆ An equivalence relation R on S satisfies the following properties
 - Reflexive:** $(x,x) \in R$
 - Symmetric:** $(x,y) \in R \Rightarrow (y,x) \in R$
 - Transitive:** $(x,y) \in R \wedge (y,z) \in R \Rightarrow (x,z) \in R$
- ◆ An equivalence relation R on S induces a partition of the elements of S into equivalence classes
- ◆ Example (connectivity relation among the vertices of a graph):
 - Let V be the set of vertices of a graph G
 - Define the relation
 - $C = \{(v,w) \in V \times V \text{ such that } G \text{ has a path from } v \text{ to } w\}$
 - Relation C is an equivalence relation
 - The equivalence classes of relation C are the vertices in each connected component of graph G

Link Relation

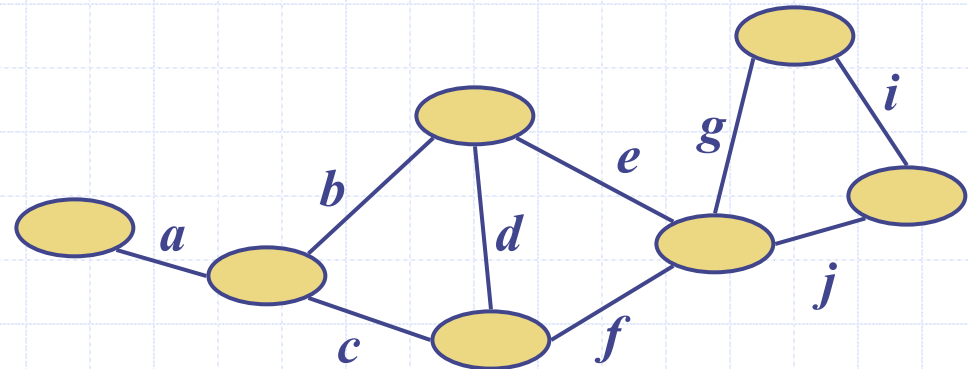
- ◆ Edges e and f of connected graph G are linked if
 - $e = f$, or
 - G has a simple cycle containing e and f

Theorem:

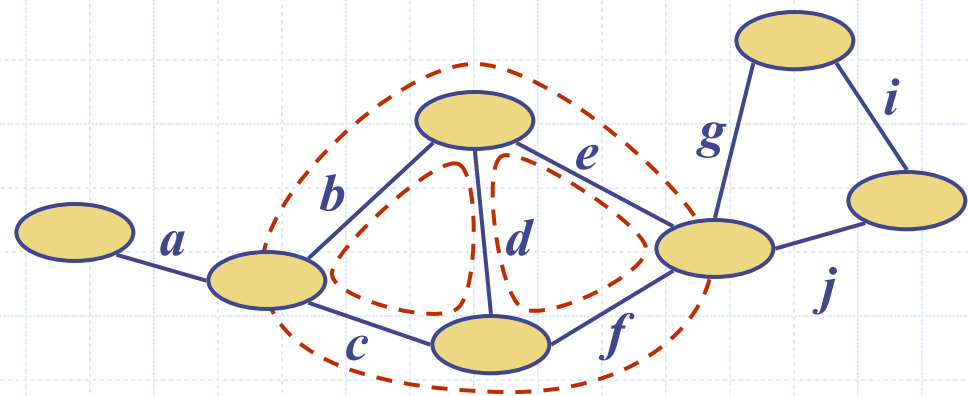
The link relation on the edges of a graph is an equivalence relation

Proof Sketch:

- The reflexive and symmetric properties follow from the definition
- For the transitive property, consider two simple cycles sharing an edge

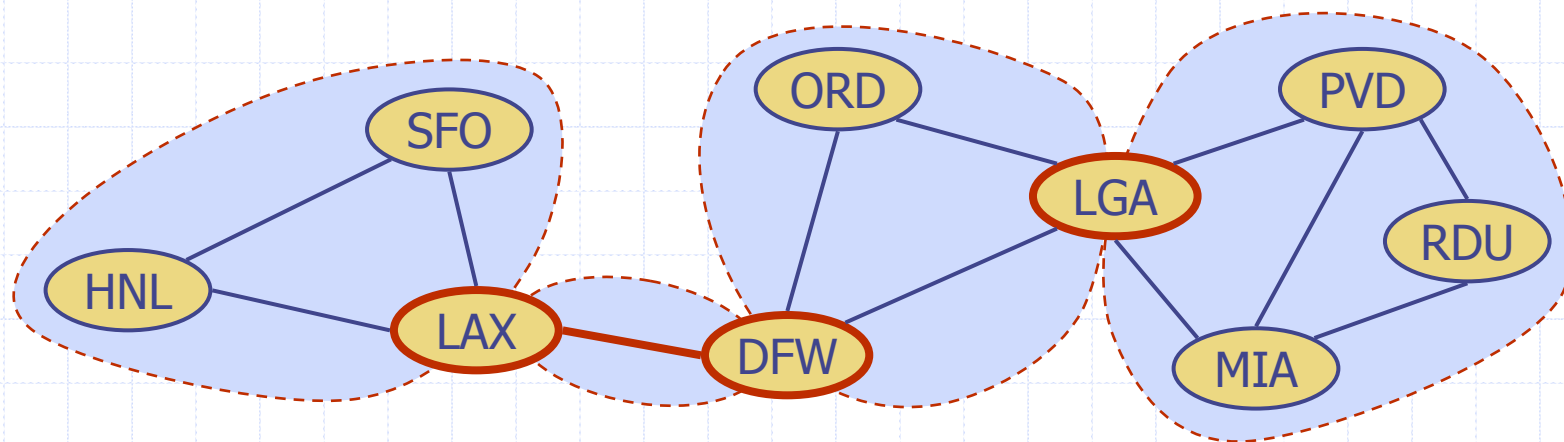


Equivalence classes of linked edges:
 $\{a\}$ $\{b, c, d, e, f\}$ $\{g, i, j\}$



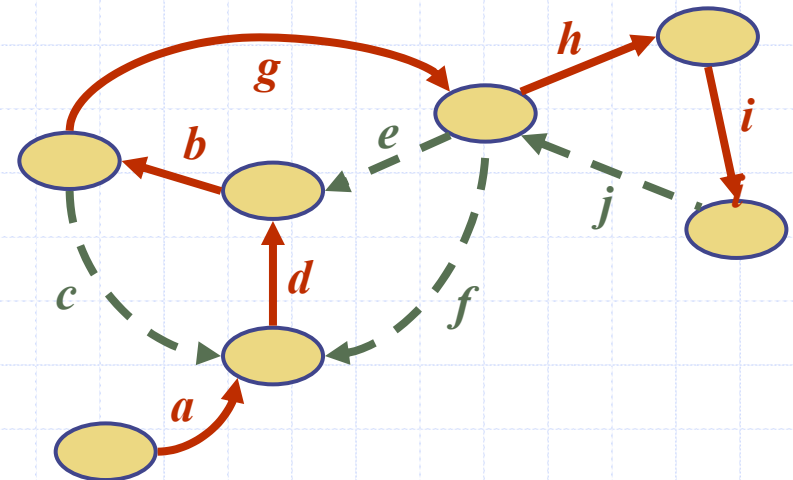
Link Components

- ◆ The link components of a connected graph G are the equivalence classes of edges with respect to the link relation
- ◆ A biconnected component of G is the subgraph of G induced by an equivalence class of linked edges
- ◆ A separation edge is a single-element equivalence class of linked edges
- ◆ A separation vertex has incident edges in at least two distinct equivalence classes of linked edge

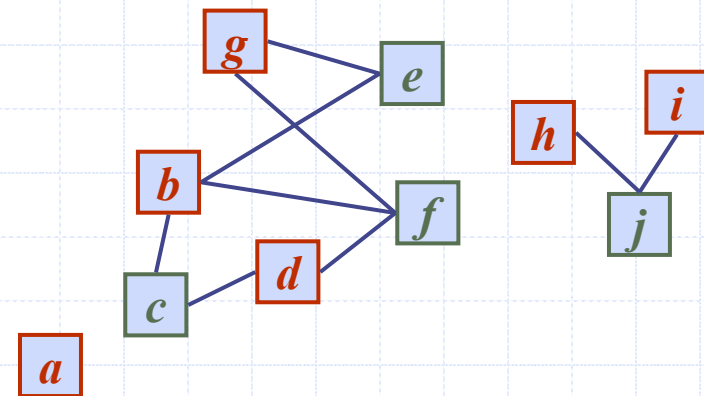


Auxiliary Graph

- ◆ Auxiliary graph B for a connected graph G
 - Associated with a DFS traversal of G
 - The vertices of B are the edges of G
 - For each back edge e of G , B has edges $(e, f_1), (e, f_2), \dots, (e, f_k)$, where f_1, f_2, \dots, f_k are the discovery edges of G that form a simple cycle with e
 - Its connected components correspond to the link components of G



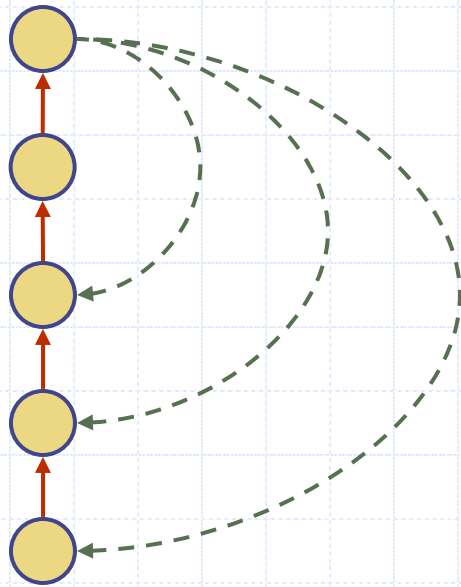
DFS on graph G



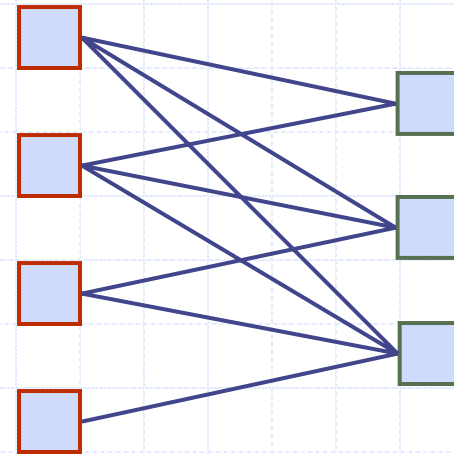
Auxiliary graph B

Auxiliary Graph (cont.)

- ◆ In the worst case, the number of edges of the auxiliary graph is proportional to nm



DFS on graph G



Auxiliary graph B

Proxy Graph

Algorithm *proxyGraph*(G)

Input connected graph G

Output proxy graph F for G

$F \leftarrow$ empty graph

DFS(G, s) { s is any vertex of G }

for all discovery edges e of G

$F.insertVertex(e)$

$setLabel(e, UNLINKED)$

for all vertices v of G in DFS visit order

for all back edges $e = (u, v)$

$F.insertVertex(e)$

repeat

$f \leftarrow$ discovery edge with dest. u

$F.insertEdge(e, f, \emptyset)$

if $f.getLabel(f) = UNLINKED$

$setLabel(f, LINKED)$

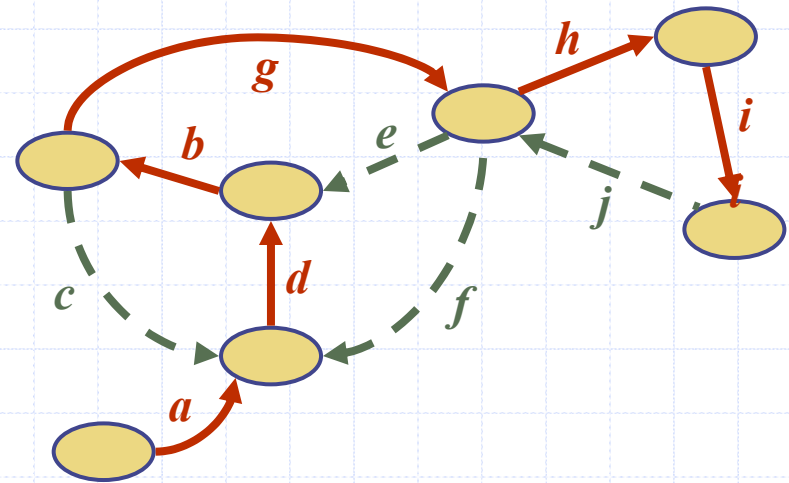
$u \leftarrow$ origin of edge f

else

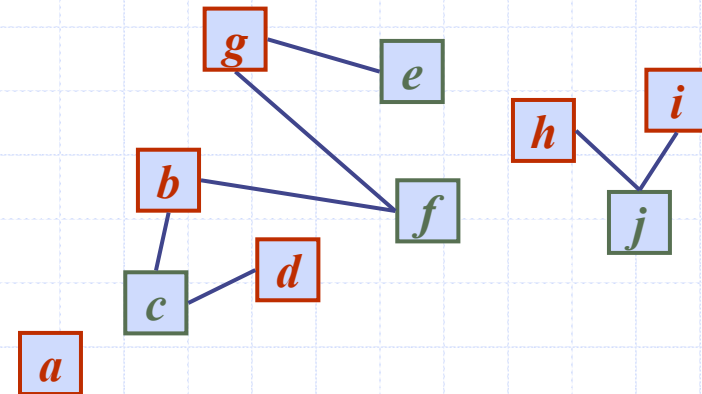
$u \leftarrow v$ { ends the loop }

until $u = v$

return F



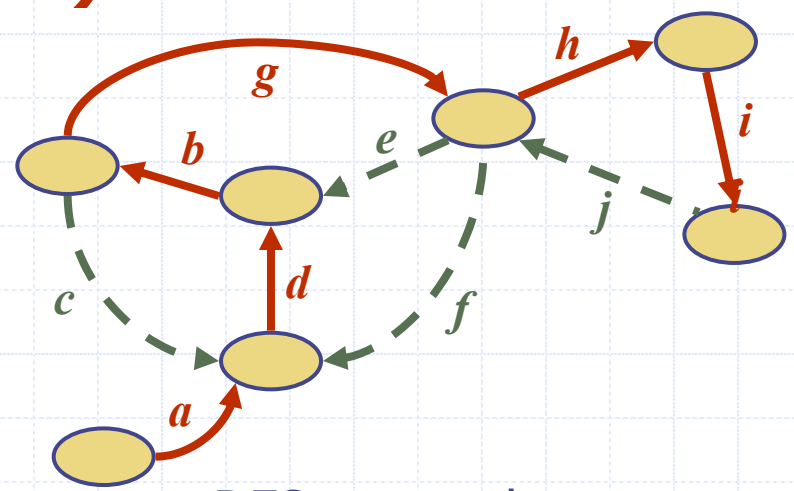
DFS on graph G



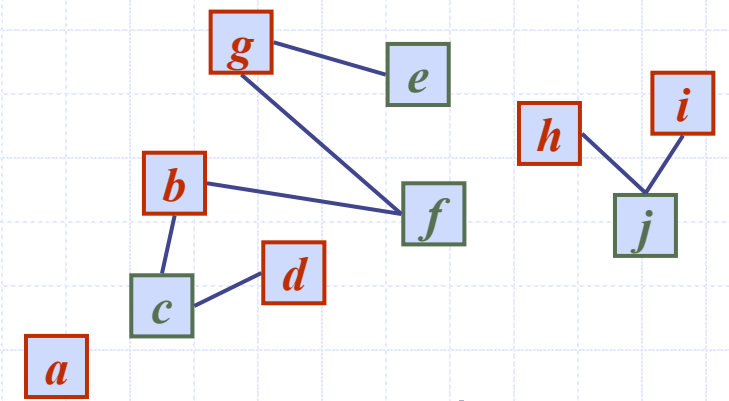
Proxy graph F

Proxy Graph (cont.)

- ◆ Proxy graph F for a connected graph G
 - Spanning forest of the auxiliary graph B
 - Has m vertices and $O(m)$ edges
 - Can be constructed in $O(n + m)$ time
 - Its connected components (trees) correspond to the link components of G
- ◆ Given a graph G with n vertices and m edges, we can compute the following in $O(n + m)$ time
 - The biconnected components of G
 - The separation vertices of G
 - The separation edges of G



DFS on graph G



Proxy graph F